



**LOYOLA COLLEGE (AUTONOMOUS), CHENNAI – 600 034**

**B.Sc. DEGREE EXAMINATION – MATHEMATICS**

FIFTH SEMESTER – NOVEMBER 2011

**MT 5407 - FORMAL LANGUAGES AND AUTOMATA**

Date : 12-11-2011  
Time : 9:00 - 12:00

Dept. No.

Max. : 100 Marks

**PART – A**

**Answer ALL questions (10 x 2 = 20)**

1. Define a finite automaton.
2. Construct the state diagram for the automaton  $M = (\{q_0, q_1, q_2\}, \{a, b\}, \delta, q_0, \{q_2\})$  where  $\delta$  is given by

$$\begin{array}{ll} \delta(q_0, a) = q_1 & \delta(q_1, a) = q_1 \\ \delta(q_2, a) = q_1 & \delta(q_0, b) = q_2 \\ \delta(q_1, b) = q_2 & \delta(q_2, b) = q_0 \end{array}$$

3. Define a non deterministic finite automaton.
4. Prove that any finite subset is regular.
5. Define context-sensitive language.
6. Write a grammar for the language  $L = \{a^n b^n / n \geq 1\}$ .
7. Define concatenation of two languages.
8. Define an  $\varepsilon$  – free homomorphism.
9. State the Chomsky Normal form for the regular expressions.
10. Define ambiguously derivable.

**PART – B**

**Answer any FIVE questions (5 x 8 = 40)**

11. Let  $M = (Q, \Sigma, \delta, q_0, F)$  be a finite automaton. Let  $R$  be a relation in  $Q$  defined by  $q_1 R q_2$  if  $\delta(q_1, a) = \delta(q_2, a)$  for all  $a$  in  $\Sigma$ . Show that  $R$  is an equivalence relation.
12. Construct a finite automaton that accepts exactly those input strings of 0's and 1's that end in 11.
13.  $M = (\{q_0, q_1, q_2, q_3\}, \{0, 1\}, \delta, q_0, \{q_3\})$  is a finite automaton and  $\delta$  is given by

$Q \setminus \Sigma$	0	1
$q_0$	$q_1$	$q_1$
$q_1$	$q_2$	$q_3$
$q_2$	$q_3$	$q_1$
$q_3$	$q_0$	$q_2$

Find (i)  $\hat{\delta}(q_0, 011010)$

(ii)  $\hat{\delta}(q_0, 1110011)$

14. Prove that union of two regular sets is also regular.

15. Let  $G = (N, T, P, S)$ ,  $N = \{S, B\}$ ,  $T = \{a, b, c\}$ .  $P$  consists of the following productions:

1.  $S \rightarrow aSBc$
2.  $S \rightarrow abc$
3.  $cB \rightarrow Bc$
4.  $bB \rightarrow bb$

Then show that  $L(G) = \{a^n b^n c^n / n \geq 1\}$  is a CSL.

16. Write a CNF grammar for the language  $L = \{wcw^R / w \in (a, b)^*\}$  and give two examples.

17. Prove that if  $L$  is a CFL generated by  $G = (N, T, P, S)$ , where  $P$  consists of rules of the form  $A \rightarrow \alpha$ ,  $A \in N$ ,  $\alpha \in (N \cup T)^*$ , then  $L$  can be generated by a CFG in which every rule is either of the form  $A \rightarrow \alpha$ ,  $A \in N$ ,  $\alpha \in (N \cup T)^+$ , or  $S \rightarrow \epsilon$ , Further  $S$  does not appear on the right side of any rule.

18. Let  $G = (N, T, P, S)$ , where  $N = \{S, A\}$ ,  $T = \{a, b\}$  and  $P$  consists of the rules

1.  $S \rightarrow aAb$
2.  $S \rightarrow abSb$
3.  $S \rightarrow a$
4.  $A \rightarrow bS$
5.  $A \rightarrow aAAb$

Find the leftmost and rightmost derivations for the string  $abab$ .

### PART – C

#### Answer any TWO question (2 x 20 = 40)

19. a) Let  $M = (\{q_0, q_1, q_2, q_3\}, \{a, b\}, \delta, q_0, \{q_1\})$  where  $\delta$  is given by

$$\delta(q_0, a) = q_1 \qquad \delta(q_0, b) = q_2$$

$$\delta(q_1, a) = q_3 \qquad \delta(q_1, b) = q_0$$

$$\delta(q_2, a) = q_2 \qquad \delta(q_2, b) = q_2$$

$$\delta(q_3, a) = q_2 \qquad \delta(q_3, b) = q_2$$

(a) Construct the state table for the given automaton  $M$ .

(b) Draw the state diagram for the given automaton  $M$ .

(c) Which of the following strings are accepted by  $M$ ?

- (i) ababa      (ii) aabba      (iii) aaaab      (iv) bbbaa

b) Construct a finite automaton  $M$  accepting  $\{ab, ba\}$ . (15+5)

20. State and prove the Pumping Lemma. (20)

21. Let  $G = (\{S, Z, A, B\}, \{a, b\}, P, S)$  where P consists of the following productions:

- |                        |                        |                        |
|------------------------|------------------------|------------------------|
| 1. $S \rightarrow aSA$ | 4. $Z \rightarrow bB$  | 7. $bB \rightarrow bb$ |
| 2. $S \rightarrow aZA$ | 5. $BA \rightarrow AB$ | 8. $bA \rightarrow ba$ |
| 3. $Z \rightarrow bZB$ | 6. $AB \rightarrow Ab$ | 9. $aA \rightarrow aa$ |

Then show that  $L(G) = \{a^n b^m a^n b^m / n, m \geq 1\}$ . (20)

22. a). Prove that the family of CFL is closed under substitution.

b). Let  $G = (N, T, P, S)$ , be any context-free grammar generating a non- empty language.

Show that there exists an equivalent grammar  $G_1$  such that for each non –terminal  $A$  of

$G_1$ , there is a derivation  $S \xRightarrow{*} \alpha_1 A \alpha_2, \alpha_1, \alpha_2 \in (NUT)^*$ . (10+10)

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